SIMULATION OPTIMIZATION USING TABU SEARCH

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ABSTRACT

Investigation of the performance and operation of complex systems in manufacturing or other environments, analytical models of these systems become very complicated. Because of the complex stochastic characteristic of the systems, simulation is used as a tool to analyze them. The trust of such simulation analysis usually is to determine the optimum combination of factors that effect the considered system performance. The purpose of this study is to use a tabu search algorithm in conjunction with a simulation model of a JIT system to find the optimum number of kanbans.

1 INTRODUCTION

When investigating the performance and operation of complex systems in manufacturing or other environments, analytical models of these systems become very complicated. Because of the complex stochastic characteristic of the systems, simulation is used as a tool to analyze them. However, the major drawback of simulation for practical applications is that is time consuming.

To overcome the limitations of simulation, metamodeling was first proposed by Blanning (1975). Major benefits of metamodeling were summarized by Madu (1990). Yu and Popplewell (1994) presented a general review of metamodels in manufacturing. Barton (1992) reviewed different methods for choosing a functional form for the metamodel relationship such as polynomials, Taguchi models and generalized linear models. Madu and Kuei (1992) employed group screening and Taguchi models in the design of experiment stage of a multi-echelon maintenance float simulation. Regression metamodels for steady-state systems are used, and they can also represent dynamic behavior in response to several unexpected real-time events in manufacturing. Lin and Cochran (1990) analyzed the dynamic performance of a hypothetical multi-station, multi-server assembly line. Lin and Chiu (1993) modeled a manufacturing cell with a fixed flow pattern considering dynamic effects of machine breakdowns and job changes. Ozdemirel et al. (1996) proposed a knowledge-based system, namely design of experiments for simulation, that assist an experienced analyst. A general review on simulation optimization was given in Tekin and Sabuncuoglu (1998).

Recently, many studies using modern heuristic techniques for simulation optimization have been encountered. Bulgak and Sanders (1988) used modified simulated annealing to optimize buffer sizes in automatic assembly systems. Grefensette (1991) considered strategy acquisition with genetic algorithms. Stuckman (1988, 1990) has discussed the use of a particular Bayesian global search algorithm for optimizing a design via simulation. Stuckman (1991) has compared three classes of global search algorithms such as genetic algorithms and Bayesian/sampling algorithms for design optimization. Haddock et al. (1992) used simulated annealing to optimize parameter levels considering the total profit of an automatic production system. Hall and Bowden (1997) presented a comparative study of direct search methods such as tabu search (TS), evolution strategies, and the Nelder Mead Simplex algorithm for simulation optimization. Lutz et al. (1995) have built a simulation model of a manufacturing process and used TS, a heuristic procedure, to optimize buffer location and storage size in this manufacturing system. Dengiz et al. (2000) used a regression metamodel to optimize batch sizes in a Printed Circuit Board assemble line.

Due to the manufacturing companies interest, researchers started investigating the JIT philosophy and also much work has been done to find number of kanbans required in a JIT system. In general, a JIT system, if implemented properly, will result in increased productivity, reduced work-in-process (WIP), and higher product quality depending on the environmental factors of the JIT system. In JIT systems both the level of WIP and order lead-time are important performance parameters. Inventory control in
a JIT system is controlled by the number of kanbans allocated. Kanban, which is a card in Japanese is used to direct materials to workstations and passes information as to what and how much to produce (Wang and Wang 1991). Kimura and Terada (1981) describe the operation of kanban systems and examine the accompanying inventory fluctuations in a JIT environment. Rees et al. (1987) studied empirical approaches for setting kanban levels dynamically. Bitran and Chang (1987) presented a mathematical formulation of the kanban determination problem. The formulation assumes planning periods of known length and finds the minimum feasible number of kanbans. Deleersnyder et al. (1989) places the kanban determination problem into the context of the overall pull system implementation problem. Monden (1981) described the model in equation 1 for setting the number of kanbans for the Toyota Motor Company.

\[ k_i = \frac{\tau_i D_i (1 + \alpha)}{n_i} \]  

(1)

where \( k_i \) is the number of kanbans for part type \( i \), \( n_i \) is the container size, \( \tau_i \) is the sum of the lead time, waiting time and kanban collecting time, and \( D_i \) is the average demand rate for part type \( i \). While the demand rate is known on average, some variability does exist due to the order sequence at final assembly and drifts in demand. Because the safety factor, \( \alpha \), in equation 1 is to handle variability, the problem is the selection of \( \alpha \). Askin et al. (1993) proposed an economic approach for selecting \( k_i \), and accordingly \( \alpha \). Their objective is to minimize the sum of inventory holding and backorder cost. They formulated a continuous time, steady-state Markov model to determine the optimal number of kanbans to use for each part type at each workcenter in a JIT system. The model selects the proper safety factor in each case.

Fukukawa and Hong (1993) proposed a mixed integer programming approach to examine many factors which play an important role in determining the number of kanbans in a JIT production system. Their objective function was to minimize inventory holding, outage and miscellaneous operating costs. Muckstadt and Tayur (1995) presented a heuristic approach to determine the optimal number of kanbans to use for each part type at each workcenter in a JIT system. The model selects the proper safety factor in each case.

The aim of this study is to find optimum number of kanbans in a JIT system using TS. The performance of TS algorithm is compared with the performance of a random search algorithm (RS) applied on the same problem. The example chosen to demonstrate this approach is a stochastic discrete-event system described by Aytug et al. (1996).

## 2 KANBAN-CONTROLLED SYSTEM

The example concerns the manufacture of two products, which are labeled Part7 and Part8. The production system includes two workcenters which are treated as black boxes (i.e. no specific flow patterns and machines are defined within the cells). The kanban controlled system considered in this study is shown in Figure 1. The products are manufactured in two adjacent workcenters. The first workcenter uses raw materials Part1, Part2, Part3, Part4 and produces two intermediate products, Part5 and Part6. The second workcenter gets the Part5 and Part6 products from workcenter 1, and produces the end products Part7 and Part8. Part5, Part6, Part7 and Part8 are produced in the manner given in Figure 1.

![Figure 1: The Kanban-Controlled Manufacturing System](image)

For each type of product at each output buffer, specific production kanbans defined as PK5, PK6, PK7, PK8, and withdrawal kanbans WK5, WK6 are used to signal requests for part transfers between the workcenters. The interarrival rate of the customer orders, order quantity and time delays in the system are given in Table 1.

<table>
<thead>
<tr>
<th>Interarrival rate of the customer orders</th>
<th>N(3, 0.6)</th>
</tr>
</thead>
<tbody>
<tr>
<td>Order quantity</td>
<td>N(1, 3)</td>
</tr>
<tr>
<td>Time Delays</td>
<td></td>
</tr>
<tr>
<td>Part5</td>
<td>N(1, 0.2)</td>
</tr>
<tr>
<td>Part6</td>
<td>N(1, 0.2)</td>
</tr>
<tr>
<td>Part7</td>
<td>N(2, 0.4)</td>
</tr>
<tr>
<td>Part8</td>
<td>N(1, 0.2)</td>
</tr>
</tbody>
</table>

Part5 = Part1 + Part2
Part6 = Part3 + Part4
Part7 = Part5 + Part6
Part8 = Part5
Problem assumptions are:

- An infinite supply of parts 1 - 4 is in the first workcenter
- Transfer time for production kanbans (PKs) is negligible
- Withdrawal kanban (WK) transfer times are also zero (the succeeding workcenter is located very close)
- Container size is one for all the parts
- The kanbans are released as soon as their containers are empty
- Customer orders (demand) are external
- Demands for Part 7 40% of the time and for Part 8 60% of the time
- Machine capacities are fixed and set at a level so that they will not generate any bottlenecks within the investigated range of the input variables

As shown in Table 1, the normal distribution is used to represent all time delays in the system to create a relatively stable environment, which is one of the main assumptions of a kanban implementation.

The objective function of this study is to minimize the possibility of backorders among workcenters and to keep the customer order cycle time at a reasonably low level. Order cycle time is defined as the difference between the time of the completion of the order and the time of the arrival of the order. Using order cycle time as a measure of performance summarizes the effect of all the internal factors in a kanban system. The order cycle times for each end product constitute the response variable values, which are obtained by a simulation model for Part 7 and Part 8 separately.

3 THE PROBLEM FORMULATION

A simulation model for the defined example was coded in PASCAL to obtain the average cycle time of Part 7 and Part 8 for each kanban combination. The problem is given in equation 2.

\[
\begin{align*}
\text{Minimize} & \quad f = \text{we} \cdot \text{tnop} + \text{kc} \cdot \text{totkan} + \text{makespan} \\
\text{Subject to} & \quad 1 \leq \text{PK5} \leq 6 \\
& \quad 1 \leq \text{PK6} \leq 3 \\
& \quad 1 \leq \text{PK7} \leq 6 \\
& \quad 1 \leq \text{PK8} \leq 6 \\
& \quad 1 \leq \text{WK5} \leq 6 \\
& \quad 1 \leq \text{WK6} \leq 3 \\
& \quad 0.4\text{cyc7} + 0.6\text{cyc8} \leq \text{maxcyc} \\
& \quad \text{totkan} = \text{PK5} + \text{PK6} + \text{PK7} + \text{PK8} + \text{WK5} + \text{WK6} \\
& \quad \text{PK5,PK6,PK7,PK8,WK5,WK6 are integers}
\end{align*}
\]

where:

- \( f \): total cost under given combination of kanbans,
- \( \text{we} \): waiting cost per minute and per order,
- \( \text{tnop} \): total number of order processed,
- \( \text{kc} \): kanban cost per minute,
- \( \text{maxcyc} \): upper limit on the average order cycle time,
- \( \text{cyc7} \): order cycle time of Part 7 (obtained from the simulation model),
- \( \text{cyc8} \): order cycle time of Part 8 (obtained from the simulation model),
- \( \text{act} \): average cycle time for both part times = 0.4\text{cyc7} + 0.6\text{cyc8}
- \( \text{makespan} \): completion time of all the orders.

Total cost \( f \) is a function of the number of kanbans only. Order cycles \( \text{cyc7} \) and \( \text{cyc8} \) are also functions of the number of kanbans only. Makespan could also be determined as a function of the number of kanbans in the system. The makespan values for different combinations were very similar with only a small standard deviation. Thus, the makespan value is computed as an average value experimentally and considered as a constant for each replication.

All definitions, input values and parameters of the kanban controlled system defined above are considered as the same of the example given in Aytug et al. (1996). Under these conditions, our simulation model was built and validated according to the simulation model results given in Aytug et al. (1996).

Examining kanbans, PK5, PK6, PK7, PK8, WK5 and WK6, for this example means that at least 11,664 kanban combination points may need to be studied if the results of the simulation model have a small standard deviation (meaning that no simulation replication is necessary). If the problem is expanded to a more complex form such as adding one additional workcenter at the end of the line, that is, an additional four kanbans (WK7, WK8, PK9, PK10) in the range of \( 1 \leq \text{WK7} \leq 3 \), \( 1 \leq \text{WK8} \leq 3 \), \( 1 \leq \text{PK9} \leq 3 \), \( 1 \leq \text{PK10} \leq 3 \), then the problem would require examining over 944784 different kanban combination points.

Due to the combinatorial nature of the number of kanbans optimization problem, the problem can become highly complex as the number of kanbans increase. It is not practical from a computational point of view to search the complete set of all combination points. The approach considered in this study avoids having to evaluate all kanban combinations by employing a TS algorithm. As mentioned previously, the developed TS algorithm employed interacts the simulation model of the JIT system.
4 TABU SEARCH

Tabu Search (TS) is a meta-heuristic that guides a local heuristic search strategy to explore the solution space beyond local optimality. The local procedure is a search that uses an operation called a move to define the neighborhood of any given solution. The neighborhood of the current solution is explored and the best solution is selected as the new current solution. This strategy allows the search to escape the neighborhood is selected, even if it is worse than the current solution. The best solution in the current solution is explored and the best solution is neighborhood of any given solution. The neighborhood of that uses an operation called a move to define the beyond local optimality. The local procedure is a search heuristic search strategy to explore the solution space Tabu Search (TS) is a meta-heuristic that guides a local optimality.

4.1 Implementation of the TS Algorithm

The notations used in the developed algorithm are introduced below:

\[ y_0: \text{ initial solution} \]
\[ y: \text{ current solution} \]
\[ y': \text{ neighbor solution} \]
\[ y_{best}' : \text{ best neighbor solution} \]
\[ y_{best}: \text{ best solution} \]
\[ M(y): \text{ a move that yields solution } y \]
\[ ttim: \text{ tabu tenure of increase moves} \]
\[ ttdm: \text{ tabu tenure of decrease moves} \]

Solutions of the problem (candidate kanban combinations) are represented by an array with six elements. This six elements include the number of kanban for each kanban type, respectively, PK5, PK6, PK7, PK8, WK5, WK6. Increasing and decreasing feasible moves were used to obtain neighbors of any solution. The possible neighbors of a sample solution \([2, 3, 4, 3, 5, 3]\) are given in Table 2.

Table 2. The Neighbors of Solution \([2, 3, 4, 3, 5, 3]\)

<table>
<thead>
<tr>
<th>([1, 3, 4, 3, 5, 3])</th>
<th>([2, 3, 3, 3, 5, 3])</th>
<th>([2, 3, 4, 3, 3])</th>
</tr>
</thead>
<tbody>
<tr>
<td>([3, 3, 4, 3, 5, 3])</td>
<td>([2, 3, 4, 3, 6, 3])</td>
<td>([2, 2, 4, 3, 5, 3])</td>
</tr>
<tr>
<td>([2, 4, 4, 3, 5, 3])</td>
<td>([2, 3, 4, 2, 5, 3])</td>
<td>([2, 3, 4, 3, 5, 2])</td>
</tr>
<tr>
<td>([2, 4, 4, 3, 5, 3])</td>
<td>([2, 3, 4, 4, 3, 5])</td>
<td>([2, 3, 4, 3, 5, 4])</td>
</tr>
</tbody>
</table>

The initial solution in the TS algorithm was randomly selected and starting from the initial solution, all possible neighbors of the current solution are examined at each iteration, because the number of neighbors is not too large. The TS algorithm calls the simulation model to compute the total cost correspond to considered neighbor solution. Then the best neighbor is selected as the new current solution, if the neighbor is not selected via a tabu move. The tabu list drives the search to different regions of the search space. After any increasing or decreasing move operation, the activated kanban are recorded on the tabu list. Two different tabu lists called \(tabu\_increase\_start\) and \(tabu\_decrease\_start\) were built to record the tabu conditions associated with the moves a selected increase or decrease move at iteration \(i\), \(tabu\_decrease\_start\) (increase move) = \(i\), or \(tabu\_increase\_start\) (decrease move) = \(i\). The tabu tenure of increasing moves \(ttim\) and tabu tenure of decreasing moves \(ttdm\) are the number of iterations that forbid a kanban number to be increased and decreased, respectively. Any tested increase or decrease move is tabu at current iteration \(j\) if \(tabu\_increase\_start(test\_increase\_move) + ttim \geq j\) or \(tabu\_decrease\_start(test\_decrease\_move) + ttdm \geq j\). This structure is also called short term memory, which is recency based (Glover 1989). An aspiration criterion was used to decide when the tabu rule can be overridden. The aspiration criterion used in this study removes the tabu condition when any tested move yields a better solution than best solution obtained so far.

The parameters values of TS algorithm were determined experimentally to be \(ttim = 3\) iterations and \(ttdm = 4\) iterations. The developed TS algorithm is terminated when a chosen maximum iteration number is reached. The steps of the TS algorithm are follows:

Algorithm:

Step 1. Choose the initial solution \(y_0\). Current solution \(y = y_0\) and best solution \(y_{best} = y_0\). \(i = 0\) and start with empty tabu lists.

Step 2. Repeat

Step 2.1. Generate the neighbors, \(y'\), for current solution \(y\) and call the simulation model to calculate the cost function \(f(y')\).

Step 2.2. Select the best neighbor \(y'_{best}\). If \(f(y'_{best}) < f(y_{best})\) then \(y_{best} = y'_{best}\) and go to step 2.4.

Step 2.3. If \((M(y'_{best})\) is tabu) and \((f(y'_{best}) > f(y_{best}))\) then \(y'_{best} = \infty\) and go to step 2.2. else current solution \(y = y'_{best}\).

Step 2.4. \(i = i + 1\). Keep \(M(y'_{best})\) on associated tabu list for associated tabu tenure.

Step 3. Until \(i \geq i_{max}\).

5 RESULTS AND ANALYSIS

In this study, the performance of TS algorithm is compared with the performance of RS algorithm. RS is the simplest
heuristic search method that just samples solution space randomly. Thus, randomly generated kanban combination is used as input value for the simulation model of kanban-controlled system to obtain objective function value. The objection function value is compared with the minimum objective function value yet encountered. If it is smaller than the best one, it is stored as the new best solution. This process is repeated until a predetermined number of solutions have been generated.

The results show that the TS algorithm and the RS algorithm find the same kanban combination [1, 1, 3, 3, 1, 1] as the best result, while searching 29 and 794 solutions, respectively. The comparison among the heuristics is presented in Table 3 on the basis of three values: coefficient of variation of cost that is found by the relevant heuristic in each replication, number of solutions searched by the relevant heuristic until finding the best solution (best kanban combination that gives minimum cost) and fraction searched %.

In order to ensure a comparable computational effort devoted to each heuristic, the stopping criterion has been defined as a number of solutions visited. This number has been set to 900 in this study. Figure 2 represents the convergence of the TS and the RS algorithms. This figure and also Table 3 showed that the TS algorithm performed better than the RS with the same number of solution visited, TS converged very quickly visiting only 0.249% of the search space. On the other hand, the nonlinear regression model used by Aytug et al. (1996) to find optimum kanban numbers for this problem fitted on the 64 simulation results ($2^6 = 64$) and they also used 6 additional simulation runs to represent validation of their regression metamodel. TS used in this study, only searched 29 points (means that 29 simulation runs) to find optimum or very close optimum solution.

Table 3: Results of the Computational Experiments

<table>
<thead>
<tr>
<th></th>
<th>TS</th>
<th>RS</th>
</tr>
</thead>
<tbody>
<tr>
<td>Coefficient of variation</td>
<td>0</td>
<td>11.8</td>
</tr>
<tr>
<td>Solution searched</td>
<td>29</td>
<td>798</td>
</tr>
<tr>
<td>Fraction searched %</td>
<td>0.249</td>
<td>6.842</td>
</tr>
</tbody>
</table>

Figure 2: Convergence of the TS and RS Algorithms

6 CONCLUSIONS

In this study, a simulation searched heuristic procedure based on TS was developed and compared with a RS algorithm applied on the same problem considering both solution quality and computational efficiency for determining the optimum number of kanbans to meet production demands in a JIT system. Results indicate that TS algorithm outperforms the RS algorithm searching only 0.249% of the solution space of this problem. The result encourages us to use TS method for simulation optimization.

REFERENCES


AUTHOR BIOGRAPHIES

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